# Imbalances in k-Colorations

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## 1. INTRODUCTION

The following problem is due to Paul Erdös [1]. Color the edges of a complete graph K on n vertices red and blue. What is the largest t such that we may always find a complete subgraph in which | red edges - # blue edges | > t?

We need a more precise and more general formulation. For any set V, define

$$\mathbf{v}^{\mathbf{k}} = \{ \mathbf{W} : \mathbf{W} \subseteq \mathbf{V}, \ |\mathbf{W}| = \mathbf{k} \}.$$
 (1)

Note that  $V^2$  is the complete graph generated by V.  $V^k$  is called the complete k-graph generated by V. The elements of  $V^k$  are called k-edges. We color the k-edges. A coloring of a set A, |A| = n, is given by a map

$$g_k : A^k \to \{+1, -1\}.$$
 (2)

The values +1, -1 may be thought of as Red and Blue. The subscript k indicates a function on k-edges and will be dropped when there is no confusion. The function  $g_k$  induces another

function, also denoted by  $g_k$ , on the subsets of A given by

$$g_{k}^{(B)} = \sum_{\substack{W \subseteq B \\ |W| = k}} g_{k}^{(W)}$$
(3)

Set

$$H_{k}(n) = \min \max |g_{k}(B)|$$

$$g_{k} \xrightarrow{B \subseteq A}$$
(4)

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where  $A = \{1, ..., n\}$  and  $g_k$  ranges over all functions satisfying (2). Clearly  $H_2(n)$  is the t required in the opening paragraph. Erdös [1] showed  $\frac{n}{4} \leq H_2(n) \leq cn^{3/2}$ . We prove

Theorem: For 
$$k \ge 1$$
, and n sufficiently large  

$$C_k n^{(k+1)/2} \le H_k(n) \le C'_k n^{(k+1)/2}$$
(5)

where the  $C_{\mu}$ ,  $C'_{\mu}$  are positive absolute constants.

2. THE PROOF

The case k = 1 is trivial,  $H_1(n) = \{\frac{n}{2}\}$ . We sketch the proof of the upper bound.

Fix  $B \subseteq A$ , |B| = b. Letting  $g_k$  be random,  $g_k(B)$  is the sum of  $\binom{b}{k}$  values  $g_k(W)$ . The values  $g_k(W)$  are +1 with probability  $\frac{1}{2}$ , -1 with probability  $\frac{1}{2}$  and independent. Thus the distribution of  $g_k(B)$  may be approximated by a normal curve of

mean 0 and 
$$\sigma = {\binom{b}{k}}^{1/2} \le n^{k/2}$$
. Thus  
Prob  $[|g_k(B)| \ge cn^{(k+1)/2}] \le e^{-c^2 n/2}$ . (6)

As there are 2<sup>n</sup> choices of B

Prob 
$$[\max_{B \subseteq A} |g_k(B)| \ge cn^{(k+1)/2}] \le 2^n e^{-c^2 n/2}$$
. (7)

For  $c = \sqrt{2 \log 2}$  the right hand side of (7) is less than unity so there does exist  $g_k$  such that max  $|g_k(B)| \leq cn^{(k+1)/2}$ . A more careful proof, using that fact that most  $|B| \sim \frac{n}{2}$  will show the upper bound of (5) for

$$C_{k}' \sim \frac{\sqrt{\log 2}}{2^{(k+1)/2} k!^{1/2}}$$
 (8)

Let us define

$$g_{k}(B_{1}^{a_{1}}B_{2}^{a_{2}}...B_{t}^{a_{t}}) = \sum g_{k}(W)$$
 (9)

where the sum is over all  $W \subseteq A$ , |W| = n,  $|W \cap B_i| = a_i$  for  $1 \le i \le t$ . This shall only be defined when the  $B_i$  are disjoint and  $\sum_{i=1}^{t} a_i = k$ .

Now we give a quick proof of the lower bound (5) for k = 2. Applying the methods of [3] we find sets  $B_1$ ,  $B_2 \subseteq A$  with

$$g_2(B_1B_2) \ge cn^{3/2}$$
 (10)

for some absolute constant c. But

$$g_2(B_1) + g_2(B_2) + g_2(B_1B_2) = g_2(B_1 \cup B_2)$$
 (11)

so

$$|g_2(V)| \ge \frac{c}{3} n^{3/2}$$
 for  $V = B_1, B_2, \text{ or } B_1 \cup B_2.$  (12)

Now we prove our theorem for all k > 2.

Lemma 1: Fix  $k \ge 2$ . Then there exists  $d_1, \ldots, d_k > 0$ ,  $t_o$ , such that for  $t > t_o$  and  $A_j$  pairwise disjoint,  $|A_j| = t$ ,  $1 \le j \le k$  we have

$$|\{(B_1, \dots, B_i) : B_j \subseteq A_j, |g_i(B_1 \dots B_i)| \ge t^{1/2}\}| \ge d_i 2^{t_i}$$
(13)

for all  $g_i$ ,  $1 \le i \le k$ .

We shall first require

Lemma 2: Fix  $c_1 > 0$ . There exists  $c_2 > 0$ ,  $t_o$ , such that  $t \ge t_o$  implies that for any choice of real  $x_j$ ,  $1 \le j \le t$ , satisfying  $|x_j| \ge 1$  for  $1 \le j \le c_1 t$ , we have

$$\left|\sum_{j\in V} x_{j}\right| > \sqrt{t} \tag{14}$$

for at least  $c_2^2$  choices of  $V \subseteq \{1, \ldots, t\}$ .

*Proof:* For  $V \subseteq \{1, \ldots, t\}$  set  $\varphi(V) = \sum_{j \in V} x_j, V_1 = V \cap \{j : 1 \\ < j < c_1 t\}, V_2 = V - V_1$ . Then  $\varphi(V) = \varphi(V_1) + \varphi(V_2)$  so (14) does not hold if  $\varphi(V_1) \in [\varphi(V_2) - \sqrt{t}, \varphi(V_2) + \sqrt{t}]$ . By a theorem of Erdös [2] for  $V_2$  fixed this holds for at most

$$\sum_{|\mathbf{r} - \frac{c_1 t}{2}| \le \sqrt{t}} ( \begin{bmatrix} c_1 t \\ 1 \end{bmatrix} ) < (1 - c_2)^2 \begin{bmatrix} c_1 t \\ 1 \end{bmatrix}$$
(15)

values of  $V_1$ , where  $c_2$  is a positive constant dependent only on  $c_1$ . Summing over all  $V_2$  yields Lemma 2. Q.E.D.

Proof of Lemma 1: We use induction on i. For i = 1 set  $x_j = g_1(\{j\})$  and apply Lemma 2. Now assume Lemma 1 holds for i - 1. Any point a  $\varepsilon$  A generates, with any  $g_i$ , a coloring  $g_{i-1}^{(a)}$  on  $A - \{a\}$ . The coloring is given by

$$g_{i-1}^{(a)}(W) = g_i(W \cup \{a\}).$$

Set

$$V = \{ ((B_1, \dots, B_{i-1}), a) : B_j \subseteq A_j, a \in B_i, g_{i-1}^{(a)}(B_1 \dots B_{i-1}) \ge t^{(i-1)/2} \}.$$

We double count

$$|v| = \sum_{a} |\{(B_1, \dots, B_{i-1}) : ((B_1, \dots, B_{i-1}), a) \in v\}|$$
 (16)

$$= \sum_{\substack{B_{1}, \dots, B_{i-1}}} |\{a : ((B_{1}, \dots, B_{i-1}), a) \in V\}|.$$
(17)

By induction the inner summation (16) is at least  $d_{i-1}^{2^{t(i-1)}}$ . Thus  $|V| \ge td_{i-1}^{2^{t(i-1)}}$ . The sum (17) has  $2^{t(i-1)}$  addends, each bounded by t. Thus for at least  $(d_{i-1}/2)2^{t(i-1)}$  choices of  $(B_1, \ldots, B_{i-1})$  we have  $|\{a : ((B_1, \ldots, B_{i-1}) | a) \in V\}| \ge d_{i-1}^{t/2}$ . Fix such a  $(B_1, \ldots, B_{i-1})$ . Set

$$x_{a} = \frac{g_{i}(B_{1} \cdots B_{i-1} \{a\})}{t^{(i-1)/2}}, \quad a \in B_{i}.$$

By assumption  $|\mathbf{x}_a| \ge 1$  for at least  $d_{i-1}t/2$  of the a. By Lemma 2 there exists  $c_2$  such that

$$|g(B_1 \cdots B_{i-1}B_i)| = |\sum_{a \in B_i} g(B_1 \cdots B_{i-1}\{a\})|$$
$$= t^{(i-1)/2} |\sum_{a \in B_i} x_a|$$

$$\geq t^{i/2}$$

for  $C_2^{2^{t}}$  choices of  $B_i$ . As this is true for at least  $(d_{i-1}^{2})^{2^{t}(i-1)}$  choices of  $(B_1, \ldots, B_{i-1}^{2})$  we may show (15) for  $d_i = d_{i-1}c_2^{2}$ , completing the induction. Q.E.D.

Now let  $g = g_k$  be any coloring on A, |A| = n. For  $t = [\frac{n}{k}]$  find disjoint  $A_1, \ldots, A_k \subseteq A$ ,  $|A_1| = t$ . From the proof of Lemma 1 we find, and fix  $B_1, \ldots, B_{k-1}$  and  $\delta > 0$  such that

$$|\{a : |g(B_1 \dots B_{k-1} \cdot \{a\})| \ge t^{(k-1)/2}\}| \ge 2\delta t.$$
 (18)

Either  $\delta t$  a's have  $g(B_1 \ \dots \ B_{k-1} \ \cdot \ \{a\}) \ge t^{(k-1)/2}$  or  $\delta t$  a's have  $g(B_1 \ \dots \ B_{k-1} \ \cdot \ \{a\}) \le -t^{(k-1)/2}$ . By symmetry (between g and -g) assume the former. Set  $B_k = \{a : g(B_1 \ \dots \ B_{k-1} \ \cdot \ \{a\}) > n^{(k-1)/2}\}$ . Then

$$g(B_{1} \cdots B_{k}) = \sum_{a \in B_{k}} g(B_{1} \cdots B_{k-1} \cdot \{a\})$$

$$\geq \delta t^{(k+1)/2}$$

$$\geq \varepsilon n^{(k+1)/2}$$
(19)

where  $\varepsilon \sim \delta/k^{(k+1)/2} > 0$ , independent of n.

To prove our result we first need a result in polynomial approximations. If G is a polynomial in, say, s variables we set |G| = the maximum absolute value of a coefficient of G and  $||G|| = \max \{G(x_1, \dots, x_s) : 0 \le x_i \le 1 \text{ for } 1 \le i \le s\}.$ 

Lemma 3: There exists  $\varepsilon = \varepsilon(s) > 0$  such that if G is a polynomial in s variables with degree at most s then

$$||G|| \ge \varepsilon |G| \tag{20}$$

*Proof:* Set  $T = \{G : |G| = 1\}$ . With the  $|\cdot|$  metric, T is compact,  $||\cdot||$  is continuous, non zero, so there exists  $\varepsilon$ , |G| = 1 => $||G|| \ge \varepsilon$ . But any  $G = |G| G_1, G_1 \varepsilon T$ , so  $||G|| = |G| ||G_1|| \ge \varepsilon |G|$ .

It should be noted that by other methods explicit bounds on  $\varepsilon(s)$  may be found.

*Proof of Theorem:* We need transfer the imbalance (19) of a product into the imbalance of a set. For  $1 \le i \le k$  let  $W_i$  range over all subsets of  $B_i$ ,  $|W_i| = [x_i|B_i|]$ , where  $0 \le x_i \le 1$  will be determined later.

$$g(W_1 \cup \ldots \cup W_k) = \sum g(W_1^{a_1} \ldots W_k^{a_k})$$
(21)

where the summation ranges over all nonnegative integers  $a_i$ ,  $\sum_{i=1}^{k} a_i = k$ . Fix the  $a_i$ . Set  $v(V_1, \dots, V_k, W_1, \dots, W_k) = 1$  if  $V_i \subseteq W_i$  for  $1 \le i \le k$ 0 otherwise.

Then from (9) the expected value

$$E[g(W_1^{a_1} \dots W_k^{a_k})] = E[\sum g(V_1 \cup \dots \cup V_k) \lor (V_1, \dots, V_k, W_1, \dots, W_k)]$$
  
(the summation over all  $V_i \subseteq B_i$ ,  $|V_i| = a_i$ )  
$$= \sum g(V_1 \cup \dots \cup V_k) E[\lor (V_1, \dots, V_k, W_1, \dots, W_k)]$$
$$= \sum g(V_1 \cup \dots \cup V_k) Prob[V_i \subseteq W_i, 1 \le i \le k | V_i \subseteq B_i, |V_i| = a_i]$$
$$= \prod_{i=1}^k x_i^{a_i} \sum g(V_1 \cup \dots \cup V_k)$$

(an approximation valid as k is fixed and n sufficiently large)

$$= g(B_1^{a_1} \dots B_k^{a_k}) \prod_{i=1}^{k} x_i^{a_i}$$
(22)

Setting

$$c_{a_1 \cdots a_k} = g(B_1^{a_1} \cdots B_k^{a_k})/n^{(k+1)/2}$$

we have, using (22) in (21)

$$E(g(W_1 \cup \cdots \cup W_n)) = n^{(k+1)/2} \sum_{a_1 \cdots a_k} a_1 \cdots a_k^{a_1} \cdots a_k^{a_k}$$

where by (19),  $|c_{1...1}| \ge \varepsilon$ . By Lemma 3 we find, and fix,  $x_1, \ldots, x_k$  so that

$$\mathbb{E}(g(\mathbb{W}_1 \cup \ldots \cup \mathbb{W}_n)) = \varepsilon_1^{(k+1)/2}$$

where  $|\varepsilon_1| \ge |c_{1...1}|\varepsilon(k) \ge \varepsilon \varepsilon(k)$  which depends only on k. By the definition of expected value we find, and fix,  $W'_1 \dots, W'_n, |W'_i| = [x_i|B_i|]$ , such that  $|g(W'_1 \cup \dots \cup W'_n)| \ge |E(g(W_1 \cup \dots \cup W_n))| \ge \varepsilon_1^{n}^{(k+1)/2}$ 

proving our theorem.

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