# On a Combinatorial Game 

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A drawing strategy is explained which applies to a wide class of combinatorial and positional games. In some settings the strategy is best possible. When applied to $n$-dimensional Tic-Tac-Toe, it improves a result of Hales and Jewett [5].

A family of sets $\left\{A_{k}\right\}$ is said to have property $B$ if there is a set $S$ which meets every $A_{k}$ and contains none of them. $m(n)$ is the smallest integer so that there is a family $\left\{A_{k}\right\}, 1 \leqslant k \leqslant m(n),\left|A_{k}\right|=n$ which does not have property $B$. It is known that $[2,3,6]$

$$
2^{n}\left(1+\frac{4}{n}\right)^{-1}<m(n)<c n^{2} 2^{n}
$$

$m(2)=3 ; m(3)=7 ; m(4)$ is not known.
Now we define a game connected with property $B$. Let $\left\{A_{k}\right\}$ be a family of sets. Let

$$
\bigcup_{k} A_{k}=S=\left\{a_{1}, \ldots, b_{1}, \ldots\right\} .
$$

The players alternately pick elements of $S$, and that player wins who first has all the elements of one of the $A_{k}$. Let $m^{*}(n)$ be the smallest integer for which there are $m^{*}(n)$ sets $\left\{A_{k}\right\},\left|A_{k}\right|=n, 1 \leqslant k \leqslant m^{*}(n)$, and so that the first player has a winning strategy.

Theorem. $\quad m^{*}(n)=2^{n-1}$.
Proof. First we exhibit $2^{n-1}$ sets of $n$ elements each so that the first player has a winning strategy. Let $G_{n}$ be such a collection. Put $G_{1}=\left\{a_{1}\right\}$
and $G_{1}{ }^{\prime}=\left\{b_{1}\right\}$. By induction define $G_{n+1}$ to be the collection of the sets $G_{n}$ and $G_{n}{ }^{\prime}$ with $a_{n+1}$ adjoined to each and define $G_{n+1}^{\prime}$ to be the collection of the sets $G_{n}$ and $G_{n}{ }^{\prime}$ with $b_{n+1}$ adjoined to each. The strategy for $G_{n}$ is clear. First pick $a_{n}$. If the second player picks $a_{i}$, you pick $b_{i}$ and vice versa. On each move the second player can block only half the remaining sets. On the $n$-th move you will complete a set.

Now we must prove that for any smaller collection of sets of $n$ elements the second player has a strategy which keeps the first player from winning. Here is the strategy: Give each element a value which is the sum of the values of each set it belongs to (which has not already been blocked by you). The value of such a set with $j$ elements remaining is $2^{n-j}$. Pick an element of largest value. To prove that the first player cannot win, we show that the sum $C$ of the values of all the sets remaining after his $i$-th move is less than $2^{n}$. So before his next move this sum $C$ is less than $2^{n}-V$, where $V$ is the sum of the values of the sets just blocked by the second player, i.e., the value of the element picked by him. Now on the first player's next move he doubles the value of each set containing the element picked, i.e., he adds the sum of their previous values $V^{\prime}$ to $C$. But clearly $V^{\prime} \leqslant V$ since $V$ was a maximum.

The same method gives the following slightly more general result: Let $\left\{A_{k}\right\}$ be a family of sets $\left|A_{i}\right|=n_{i}$ for which

$$
\sum_{i} \frac{1}{2^{n_{i}}}<1 .
$$

Then the next player has a strategy which forces a draw. On the other hand if integers $n_{i}$ are given for which

$$
\sum_{i} \frac{1}{2^{n_{i}}} \geqslant \frac{1}{2}
$$

we can find sets $A_{i},\left|A_{i}\right|=n_{i}$, so that the first player has winning strategy. Clearly without loss we may assume the sum equals $1 / 2$. After putting $a_{1}$ in each set it is again clear that we may put $a_{2}$ in "half" the sets and $b_{2}$ in the others, i.e., $a_{2}$ and $b_{2}$ have equal value. Now we may again split the collection in which $a_{2}$ occurs into two equal parts, and so on.

Let us now denote by $m_{1}{ }^{*}(n)$ the smallest integer for which there are $m_{1}{ }^{*}(n)$ sets $\left\{A_{k}\right\}, 1 \leqslant k \leqslant m_{1}{ }^{*}(n),\left|A_{k}\right|=n,\left|A_{i} \cap A_{j}\right| \leqslant 1,1 \leqslant i<j \leqslant$ $m_{1}{ }^{*}(n)$, and so that the first player has a winning strategy. We have no satisfactory estimate for $m_{1}{ }^{*}(n) . m_{1}{ }^{*}(3)=6$, the sets being $123,145,167$, $189,246,579$. Probably $m_{1}{ }^{*}(n)$ is considerably smaller than $m^{*}(n)$.

Hales and Jewett [5] investigated $n$-dimensional Tick-Tack-Toe in a hypercube of side $k$. They proved that, if $k \geqslant 3^{n}-1$ ( $k$ odd) or $k \geqslant 2^{n+1}-2$ ( $k$ even), the second player can force a tie, but for each $k$ there is an $n_{k}$ so that for $n \geqslant n_{k}$ the first player can win.

In the $n$-dimensional hypercube of side $k$ there are $\frac{1}{2}\left\{(k+2)^{n}-k^{n}\right\}$ sets which form winning lines. Thus our theorem immediately implies that the second player can force a draw if $k>c n \log n$. This result still falls short of their conjecture that the second player can force a draw if

$$
k>2\left(2^{1 / n}-1\right)^{-1} \approx \frac{2 n}{\log 2}-1
$$

It is well known $[4,1]$ that, if we color the edges of a complete graph of $n$ vertices by two colors, there always is a complete subgraph of

$$
\left[\frac{\log n}{2 \log 2}\right]
$$

vertices all of whose edges have the same color, but there does not have to be such a graph of

$$
\left[\frac{2 \log n}{\log 2}\right]
$$

vertices.
Now following Simmons we define a game called the Ramsey game connected with this property. The players alternately choose edges. That player wins who first gets all the edges of a complete graph of $k$ vertices. Ramsey's theorem implies that the game is a win for the first player for

$$
k \leqslant\left[\frac{\log n}{2 \log 2}\right]
$$

Our theorem implies that the game is a draw if

$$
2^{l}>\binom{n}{k}, \quad \text { where } \quad l=\binom{k}{2}-1
$$

i.e., it is a draw if

$$
k \geqslant(1+o(1)) \frac{2 \log n}{\log 2}
$$

We did not investigate the Ramsey game for triples since we did not succeed in getting any satisfactory result.

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