# ON GRAPHS WHICH CONTAIN ALL SPARSE GRAPHS 

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Dedicated to Professor A. Kotzig on the occasion of his sixtieth birthday

## 1. Introduction

Let $\mathscr{H}_{n}$ denote the class of all graphs with $n$ edges and denote by $s(n)$ the minimum number of edges a graph $G$ can have which contains all $H \in \mathscr{H}_{n}$ as subgraphs. In this paper we establish the following bounds on $s(n)$ :

## Theorem 1.

$$
\frac{c n^{2}}{\log ^{2} n}<s(n)<\frac{c^{\prime} n^{2} \log \log n}{\log n}
$$

for $n$ sufficiently large and $c, c^{\prime}$ some constants.
We also consider the problem of determining the minimum number of edges, denoted by $s^{\prime}(n)$, a graph can have which contains every planar graph on $n$ edges as a subgraph. We prove:

Theorem 2. $s^{\prime}(n)<c n^{3 / 2}$.
In $[1,2,3]$, two of the authors investigated the problem of determining the minimum number of edges a graph or a tree could have which contains all trees on $n$ edges as subgraphs. For a brief survey on these 'universal' graphs the reader is referred to [4].

## 2. A lower bound for $s(n)$

Let $G$ be a graph which contains all graphs on $n$ edges. Suppose $G$ has $t$ edges. Thus $G$ contains at most $\left(\begin{array}{l}\binom{n}{n}\end{array}\right)$ different subgraphs on $n$ edges.
On the other hand, $G$ contains all graphs on $n$ edges and $\lfloor n / \log n\rfloor$ vertices where $\lfloor x\rfloor$ denotes the largest integer less than or equal to $x$. There are at least

$$
\left(\left(\begin{array}{c}
\lfloor n / \log n\rfloor \\
2 \\
n
\end{array}\right)\right) \cdot \frac{1}{\lfloor n / \log n\rfloor!}
$$

different graphs with $n$ edges and $\lfloor n / \log n\rfloor$ vertices (see [5]). Therefore we have

$$
\binom{t}{n} \geqslant\left(\left(\begin{array}{c}
\lfloor n / \log n\rfloor \\
2 \\
n
\end{array}\right)\right) \cdot \frac{1}{\lfloor n / \log n\rfloor!} .
$$

By a straightforward calculation, this implies

$$
t \geqslant c n^{2} / \log ^{2} n
$$

for some constant $c$.
Hence we have shown $s(n)>c n^{2} / \log ^{2} n$.

## 3. An upper bound for $s(n)$

We will prove (by the probability method) that there exists a graph with $c n^{2} \log \log n / \log n$ edges $^{1}$ that contains all graphs with at most $n$ edges. The existence of such a graph will follow from the following sequence of observations.

Claim 1. Given positive integers $a$ and $b$ where $a<b<a \log a$ and $\log \log \log a \geqslant 1$, there is a bipartite graph $H$ with vertex set $A \cup B$ where $|A|=a$ and $|B|=b$, which satisfies the following conditions:
(i) $H$ has no more than abp edges where $p=\log \log a / \log a$;
(ii) For any $k$ disjoint subsets of $B$, say, $S_{1}, \ldots, S_{k}$, each with cardinality at most $p^{-1}$, and $2 \mathrm{kp}^{-2}<a$, we have

$$
\left|\bigcup_{i=1}^{k} N\left(S_{i}\right)\right| \geqslant k p^{-2}
$$

where

$$
N\left(S_{i}\right) \equiv\left\{v \in A: v \text { is adjacent to all vertices in } S_{i}\right\} .
$$

Proof. We consider the set of all bipartite graphs on $a$ and $b$ vertices with $a b p$ edges. For a set $S_{i} \subset B,\left|s_{i}\right|<d=p^{-1}$, the probability of a vertex $v$ in $A$ being in $N\left(S_{i}\right)$, is at least $p^{d}$. Therefore the probability of $v$ not being in any $N\left(S_{i}\right)$ is at most $\left(1-p^{d}\right)^{k}$. The

[^0]probability that there are $a-k d^{2}$ vertices in $A$ not in any $N\left(S_{i}\right)$ is at most
$$
\binom{a}{k d^{2}}\left(1-p^{d}\right)^{k\left(a-k d^{2}\right)} \leqslant 2^{a} \mathrm{e}^{-p^{d k a / 2}}
$$

Since there are at most $b^{d k}$ choices for $S_{i}, 1 \leqslant i \leqslant k$, the probability for a bipartite graph to be 'bad' is at most

$$
\begin{aligned}
b^{d \mathrm{k}} \cdot 2^{a} \mathrm{e}^{-p^{d} k a / 2} & <(a \log a)^{p^{-1} \mathrm{k}} \cdot 2^{a} \mathrm{e}^{-p^{d} k a / 2} \\
& <(a \log a)^{a \log \log a / \log a 2^{a} \mathrm{e}^{-a^{2} p^{d-2} / 4}<1^{\prime}}
\end{aligned}
$$

Therefore the required bipartite graph exists as claimed.

Claim 2. Given positive integers $a$ and $b$ where $a<b<a \log a$ and $\log \log \log a \geqslant 1$, there is $a$ bipartite graph $H$ with vertex set $A \cup B$ where $|A|=a$ and $|B|=b$ satisfying the following conditions:
(i) $H$ has no more than abp edges where $p=\log \log a / \log a$.
(ii) Let $H^{\prime}$ be a bipartite graph with vertex set $X \cup Y$ where $|X| \leqslant \frac{1}{2} a,|Y|=b$ and maximum degree $p^{-1}$. Then $H^{\prime}$ can be embedded in $H$ in the strong sense, i.e. any one-to-one map $\lambda: Y \rightarrow B$ can be extended to $\bar{\lambda}: X \cup Y \rightarrow A \cup B$ such that $\bar{\lambda}(u)$ is adjacent to $\bar{\lambda}(v)$ in $H$ if $u$ is adjacent to $v$ in $H^{\prime}$.

Proof. We take $H$ to be the graph in Claim 1. The mapping $\lambda$ will be extended to $\bar{\lambda}: X \cup Y \rightarrow A \cup B$ in the following way:

For a vertex $x$ in $X$, we define

$$
\begin{aligned}
& S(x)=\{b \in B: b \\
&=\lambda(y) \text { and } y \text { is adjacent to } x\}, \\
& M(x)=N(S(x))=\{v \in A: v \text { is adjacent to all vertices in } S(x)\} .
\end{aligned}
$$

The existence of $\bar{\lambda}$ is equivalent to a system of distinct representatives for $\{M(x)\}_{x \in X}$.
It suffices to show that for any set $X^{\prime} \subseteq X$ we have

$$
\left|\bigcup_{x \in X^{\prime}} M(x)\right| \geqslant\left|X^{\prime}\right| .
$$

This is clearly true for $\left|X^{\prime}\right| \leqslant(\log a / \log \log a)^{2}$ by property (ii) of $H$.
Now suppose $\left|X^{\prime}\right|>(\log a / \log \log a)^{2}$. Since $H^{\prime}$ is of bounded degree $d=$ $\log a / \log \log a$, for each $x$ there are at most $d^{2}$ vertices $x^{\prime}$ in $X$ with $S(x) \cap S\left(x^{\prime}\right) \neq \emptyset$. Thus there is a subset $X^{\prime \prime}$ of $X$ where $\left|X^{\prime \prime}\right| \geqslant\left|X^{\prime}\right| / d^{2}$ such that all $S(x), x \in X^{\prime \prime}$, are mutually disjoint. Therefore,

$$
\left|\bigcup_{x \in X^{\prime}} M(x)\right| \geqslant\left|\bigcup_{x \in \mathcal{X}^{\prime \prime}} M(x)\right| \geqslant \frac{\left|X^{\prime}\right| p^{-2}}{d^{2}} \geqslant\left|X^{\prime}\right| .
$$

This completes the proof of Claim 2.

Claim 3. There exists a graph $\bar{H}$ with $4 n^{2} \log \log n / \log n$ edges which contains all graphs with $n$ vertices and degree at most $\log n / \log \log n=d$.

Proof. We will construct a $d$-partite graph $\bar{H}$ as follows:
(i) $\bar{H}$ has vertex set $A_{1} \cup A_{2} \cup \cdots \cup A_{d+1}$ with $\left|A_{i}\right|=2 n / d$ for each $i$;
(ii) For each $i$, no $u, v \in A_{i}$ are adjacent;
(iii) The edges between $A_{i}$ and $A_{1} \cup A_{2} \cup \cdots \cup A_{i-1}$ form a graph described in Claim 2.

It can be easily seen that $\tilde{H}$ has at most $4 n^{2} \log \log n / \log n$ edges. It suffices to prove that any graph $G$ with degree $d$ can be embedded in $\bar{H}$. A nice result of Hajnal and Szemerédi [6] states that any graph with degree at most $d$ can be colored by $d+1$ colors in such a way that the sizes of the color classes differ by at most 1 . Suppose $G$ has color classes $C_{1}, \ldots, C_{d+1}$. We will then embed $C_{1}$ into $A_{1}, C_{2}$ into $A_{2}$, and so on, as guaranteed by Claim 2 .

Claim 4. There exists a graph $F(n)$ with $C n^{2} \log \log n \log n$ edges which contains all graphs on $n$ edges where $C$ is an absolute constant.

Proof. We will construct the graph $F(n)$ as follows:
(i) The vertex set is the disjoint union of $A$ and $B$ where $|A|=2 n \log \log n / \log n$ and $|B|=2 n$.
(ii) Every vertex $v$ in $A$ is adjacent to all vertices in $V(F(n))-\{v\}$.
(iii) The subgraph of $F(n)$ induced by $B$ is the graph, as described in Claim 3, which has $4 n^{2} \log \log n / \log n$ edges and contains all graphs with $2 n$ vertices and degree at most $d$.

It is easy to see that $F(n)$ has at most $10 n^{2} \log \log n / n^{2}$ edges. Let $G$ be an arbitrary graph on $n$ edges. $G$ has at most $2 n \log \log n / \log n$ vertices with degree more than $\log n / \log \log n$. These vertices will be embedded in $A$. The remaining part of the graph will then be embedded in $B$ as guaranteed by Claim 3.

This completes the proof of Claim 4.
Remark. If instead of using the result of Hajnal and Szemerédi, we use the simple fact that a graph on $n$ vertices and maximum degree $d$ can be $2(d+1)$ colored so that each color class has size at most $n / d$, then the resulting bound will differ from the one presented by a constant factor.

## 4. Universal graphs for planar graphs

We will use the following theorem to give an upper bound of $n^{3 / 2}$ for the universal graphs which contain all planar graphs on $n$ edges.

Separator Theorem (Lipton and Tarjan [6]). Let G be any planar graph with $n$ vertices. The vertices of $G$ can be partitioned into three sets, $A, B, C$ such that no edge joins a vertex in $B$ with a vertex in $C$, neither $B$ and $C$ contain more than $n / 2$ vertices, and $A$ contains no more than $2 \sqrt{2 n} /(1-\sqrt{2 / 3})$ vertices.

Let $G(m)$ denote the graph constructed as shown in Fig. 1.
The vertices of $G(n)$ can be partitioned into three parts, $X, Y$ and $Z$ where $|X|=$


Fig. 1.
$\left.2 \sqrt{2 n} /(1-\sqrt{2 / 3})=c_{1} \sqrt{n},|Y|=\mid V(G(\mid n / 2\rfloor)\right) \mid$ and $\left.|Z|=\mid V(G(\mid n / 2\rfloor)\right) \mid$. Any vertex in $X$ is adjacent to any vertex in $G(n)$ except itself. The induced subgraph on $Y$ is $G(\lfloor n / 2\rfloor)$ and the induced subgraph on $Z$ is $G(\lfloor n / 2\rfloor)$.

It is rather straightforward to see that any planar graph with $n$ vertices can be embedded in $G(n)$ since we can partition any planar graph into three parts, $A, B$ and $C$ as described in the Separator Theorem, and we can embed $A$ in $X, B$ in $Y$ and $C$ in $Z$.

We also note that $G(n)$ has fewer than $c_{2} n$ vertices since

$$
|V(G(n))|<2|V(G(n / 2))|+c_{1} \sqrt{n}
$$

and we can prove by induction on $n$ that

$$
|V(G(n))| \leqslant \frac{c_{1} \sqrt{2}}{\sqrt{2}-1} n\left(1-\frac{1}{\sqrt{2 n}}\right) .
$$

Now, by the construction of $G(n)$, we know that

$$
|E(G(n))| \leqslant|V(G(n))| \cdot c_{1} \sqrt{n}+2|E(G(n / 2))| .
$$

It follows by induction that $G(n)$ has fewer than $c n^{3 / 2}$ edges where $c=c_{1}^{2} \sqrt{2} /(\sqrt{2}-1)=$ 19.7607. . . Therefore we have

$$
s^{\prime}(n)<c n^{3 / 2}
$$

and Theorem 2 is proved.
We note that the obvious lower bound for $s^{\prime}(n)$ is $\frac{1}{2} n \log n$ which is the lower bound for the number of edges in graphs which contains all trees on $n$ edges (see [2]). At present we do not know any better lower bound than $\mathrm{cn} \log n$.

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[^0]:    ${ }^{1}$ Strictly speaking, we should use $3 n[\log \log n / \log n]$ or $[3 n \log \log n / \log n]$ since $|A|$ is an integer. However, we will usually not bother with this type of detail since it has no significant effect on the arguments or results.

