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RAMSEY NUMBERS FOR THE PAIR SPARSE GRAPH-PATH OR CYCLE

BY

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ABSTRACT. Let G be a connected graph on n vertices with no more than $n(1 + \epsilon)$ edges, and P_k or C_k a path or cycle with k vertices. In this paper we will show that if n is sufficiently large and ϵ is sufficiently small then for k odd

 $r(G, C_k) = 2n - 1.$

Also, for k > 2,

$$r(G, P_k) = \max\{n + [k/2] - 1, n + k - 2 - \alpha' - \delta\},\$$

where α' is the independence number of an appropriate subgraph of G and δ is 0 or 1 depending upon n, k and α' .

Introduction. Let G and H be simple graphs. The Ramsey number r(G, H) is the smallest integer n such that for each graph F on n vertices, either G is a subgraph of F or H is a subgraph of \overline{F} , the complement of F. Calculation of r(G, H) for particular pairs of graphs G and H has received considerable attention, and a survey of such results can be found in [2].

Chvátal [5] proved that if T_n is a tree on *n* vertices and K_m is a complete graph on *m* vertices, then $r(T_n, K_m) = (n - 1)(m - 1) + 1$. In [4] it was shown that if T_n is replaced by a sparse connected graph G_n on *n* vertices the Ramsey number remains the same (i.e. $r(G_n, K_m) = (n - 1)(m - 1) + 1$). For m = 3 Chvátal's theorem implies $r(T_n, K_3) = 2n - 1$. In this paper we will show that if T_n is replaced by any sparse connected graph G on *n* vertices and K_3 is replaced by an odd cycle C_k , then for appropriate *n* the Ramsey number is unchanged. In particular we will prove the following.

THEOREM. If G is a connected graph on n vertices and no more than $n(1 + \epsilon)$ edges, then

$$r(G, C_{\nu}) = 2n - 1$$

for n sufficiently large, ε sufficiently small (both depending upon k) and k odd.

This theorem falls into a larger category of results considered by Burr [3]. A graph G on n vertices is H-good if $r(G, H) = (n - 1)(\chi(H) - 1) + t(H)$, where $\chi(H)$ is the chromatic number of H and t(H) is the minimum number of vertices in any color class of any $\chi(H)$ vertex coloring of H. The above theorem states that G is C_k -good.

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Gerencsér and Gyárfas [7] proved that $r(P_n, P_m) = n + [m/2] - 1$ for $n \ge m$, where P_k denotes a path with k vertices. Also, Parsons [10] verified for appropriately large n that $r(S_n, P_m) = n + m - 2 - \delta$ (S_n is a star on n vertices), where $\delta = 0$ if m - 1 divides n - 2 and $\delta = 1$ otherwise. For large n we will generalize these two results by replacing the star and the large path by a sparse graph. The following theorem will be proved.

THEOREM. Let G be a connected graph with n vertices and no more than $n(1 + \epsilon)$ edges. Then for $k \ge 2$, n sufficiently large, and ϵ sufficiently small (both depending upon k),

 $r(G, P_k) = \max\{n + \lceil k/2 \rceil - 1, n + k - 2 - \alpha' - \delta\},\$

where α' is the independence number of an appropriate subgraph of G, and $\delta = 0$ if k - 1 divides $n - 2 - \alpha'$ and $\delta = 1$ otherwise.

We now mention some notation that will be used in the paper. Notation not specifically mentioned will follow [9]. The vertex set and edge set of a graph G will be denoted by V(G) and E(G) respectively. The edge with endvertices u and v will be written uv. A path with k vertices and endvertices u and v will be denoted by $P_k(u, v)$. If the vertices of the path are $\{v_1, v_2, \ldots, v_k\}$ then this will be expressed by (v_1, v_2, \ldots, v_k) . The corresponding cycle will be written $(v_1, v_2, \ldots, v_k, v_1)$. A path P which is a subgraph of a graph G is a suspended path of G if each vertex of P, except for its endvertices, has degree 2 in G. If a graph H is isomorphic to a subgraph of G, this will many times just be expressed as H is a subgraph of G. The graph G - H will be the subgraph of G induced by the vertices of G not in H. The graphical parameters α , δ and Δ will denote the independence number, minimum degree, and maximum degree of a graph respectively. Also as usual, $\{x\}$ and [x] will denote the least integer greater than or equal to x and the greatest integer less than or equal to x respectively.

Main theorems. There are several facts that are used in the proofs of both of the main theorems as well as in the proofs of other results. Some of these common facts are included in the next lemmas.

LEMMA 1. Let u and v be distinct vertices of a graph G on s + t vertices $(t \ge s \ge 2)$. Assume G contains a path P which is a $P_t(u, v)$ but G contains no $P_{t+1}(u, v)$.

(i) If $t \ge 2s$, then G contains a $P_{2(s/2)+1}$.

(ii) If $t \ge 5s - 1$, then \overline{G} contains a P_{2s+1} .

(iii) If $t \ge 5s - 1$, then \overline{G} contains a P_{2s-1} between any pair of vertices not in P or \overline{G} contains a K_{2s+1} .

PROOF. Denote the path P by (x_1, x_2, \ldots, x_i) and let S be the set of vertices of G not in P. Since the path P cannot be lengthened, no vertex of S is adjacent in G to two consecutive vertices of P. Thus for any pair $\{x_i, x_{i+1}\}$ of consecutive vertices of P, two of any three vertices of S are commonly adjacent in \overline{G} to either x_i or x_{i+1} .

SPARSE GRAPH-PATH OR CYCLE

(i) Since $t \ge 2s$ repeated application of the above fact implies that there exist two vertex disjoint paths Q_1 and Q_2 satisfying the following.

(a) Their first and last vertices are in S.

(b) Their vertices alternate between S and P.

(c) $S \subset V(Q_1) \cup V(Q_2)$.

With no loss of generality we can assume that Q_1 contains at least $\{s/2\}$ vertices of S. If Q_1 contains at least $\{s/2\} + 1$ vertices, then Q_1 is the desired path. If Q_1 contains precisely $\{s/2\}$ vertices of S (and $\{s/2\} - 1$ vertices of P), then Q_1 can be extended to a $P_{2(s/2)+1}$ by adjoining a vertex from P to each end of Q_1 . This is possible since each vertex of S is adjacent in \overline{G} to at least s vertices.

(ii) and (iii). If a vertex w of S is adjacent in G to x_i and x_j of P (i < j < t), then $x_{i+1}x_{j+1}$ is in \overline{G} . Otherwise, $(x_1, \ldots, x_i, w, x_j, x_{j-1}, \ldots, x_{i+1}, x_{j+1}, \ldots, x_i)$ is a $P_{i+1}(u, v)$. Therefore if a vertex w of S is adjacent in G to 2s + 1 vertices of P, the 2s vertices which are successors (the last vertex may not have a successor) of these vertices along P together with w form a K_{2s+1} in \overline{G} . We can thus assume that each vertex of S is adjacent in G to at most 2s vertices of P. This implies that each pair of vertices of S are commonly adjacent in \overline{G} to at least s - 1 vertices of P. Hence for any pair of vertices w_1 and w_2 of S there exists a $P_{2s-1}(w_1, w_2)$ in \overline{G} such that the vertices alternate between S and P. This path can be extended to a P_{2s+1} since each of w_1 and w_2 are adjacent in \overline{G} to at least 3s - 1 vertices of P.

LEMMA 2. Let G be a graph on n vertices and n + t edges. If G has no isolated vertices and no suspended path with more than s vertices, then G has at least $\{n/(2s) - 3t/2\}$ vertices of degree 1.

PROOF. Let H be the graph (possibly multigraph) obtained from G by replacing each suspended path of G by an edge. Thus if H has m vertices, then H has m + tedges and every vertex has degree 1 or at least 3. If x is the number of vertices in G (and in H) of degree 1, then

 $x + 3(m - x) \le 2(m + t).$

Since G has no suspended path with more than s vertices, G can be obtained from H by replacing each edge of H by a suspended of appropriate length less than or equal to s. Thus

$$n \leq s(m+t).$$

Therefore the two displayed equations imply

$$x \ge n/2s - 3t/2$$

which completes the proof.

LEMMA 3. Let $m \ge n$, $n_1 \ge n_2 \ge \cdots \ge n_l \ge 1$, $k \ge 3$ and $l \ge \{k/2\} - 1$ be positive integers. Let H be a graph on $n - \sum_{i=1}^{l} n_i$ vertices with l distinct vertices designated v_1, v_2, \ldots, v_l , and G a graph on n vertices obtained from H by adjoining for each i $(1 \le i \le l)$, a star with n_i edges and center at v_i . If

$$r(H, P_k) \leq m - \left(\sum_{i=1}^{\lfloor k/2 \rfloor - 1} n_i\right)(k-1),$$

then

$$r(G, P_k) \le m + [k/2] - 1.$$

PROOF. Let F be a graph on $m + \lfloor k/2 \rfloor - 1$ vertices. Assume that F does not have G as a subgraph and \overline{F} does not contain a P_k . We will show that this leads to a contradiction.

Select a path of maximal length in \overline{F} . By assumption this path has at most k - 1 vertices and the endvertices of the path are adjacent in F to all vertices not on the path. Delete the vertices of this path and then repeat this procedure a total of $(\sum_{i=1}^{\lfloor k/2 \rfloor - 1} n_i) - 1$ times. This leaves a graph F' with at least $m - (\sum_{i=1}^{\lfloor k/2 \rfloor - 1} n_i) \cdot (k - 1)$ vertices. Thus F' contains H as a subgraph. Let A be the set of endvertices of the paths deleted. Hence $|A| \ge \sum_{i=1}^{\lfloor k/2 \rfloor - 1} n_i$.

The graph F would contain G as a subgraph if there were disjoint sets of vertices S_1, S_2, \ldots, S_i ($|S_i| = n_i$) in F but not in H such that each vertex of S_i is adjacent in F to the vertex v_i for $1 \le i \le l$. Since F does not contain G as a subgraph, Hall's theorem [8] implies that there exists a set $L = \{v_{i_1}, v_{i_2}, \ldots, v_{i_l}\}$ such that

$$|N \cap V(F-H)| < \sum_{j=1}^{r} n_{i_j},$$

where N is the set of vertices which is adjacent in F to at least one of v_{ij} , $1 \le j \le r$. Clearly then, $|N \cup V(H)| < |V(H)| + \sum_{j=1}^{r} n_{ij} \le n$. Therefore \overline{F} contains a complete bipartite graph $K_{r,lk/2l}$. Also each vertex of A is adjacent in F to each v_i , $1 \le i \le l$. Since $|A| \ge \sum_{i=1}^{\lfloor k/2 \rfloor} n_i$ and $A \cap V(H) = \emptyset$, $r \ge \lfloor k/2 \rfloor$. This implies that \overline{F} contains a P_k , a contradiction.

The next lemma will be needed in the inductive step of the proof of Proposition 5. The graph G + e is a graph obtained from the graph G by adding the edge e between two nonadjacent vertices of G.

LEMMA 4. For any graph G and for $k \ge 2$,

$$r(G+e, P_k) \leq r(G, P_k) + k - 1.$$

PROOF. Let F be a graph on $r(G, P_k) + k - 1$ vertices and assume that \overline{F} does not contain a P_k . Select a path of maximal length in \overline{F} and let w be an endvertex of this path. Since \overline{F} contains no P_k , the set S consisting of w and the vertices adjacent to w in \overline{F} has at most k - 1 vertices. Therefore there is a copy of G in F which is vertex-disjoint from S. Let u and v be the endvertices in this copy of G of the edge e. Since w is adjacent in F to each vertex of this copy of G, replacing u (or v) by wgives a copy of G + e.

Let G be a connected graph on n vertices and no more than $[n(1 + \varepsilon)]$ edges. If n is sufficiently large and ε sufficiently small, then G has either a long suspended path or a large number of vertices of degree 1. This fact will be used in the proof of the following proposition. The various cases appearing in the proof are also a consequence of this fact.

PROPOSITION 5. For integers $n, k \ge 1$, let G be a connected graph with n vertices and no more than [n(1 + 1/18k)] edges. Then $r(G, P_k) \le n + 13k^5$.

SPARSE GRAPH-PATH OR CYCLE

PROOF. The result is trivial for k = 1 or 2, so assume $k \ge 3$. The proof will be by induction on *n*. Since $r(K_n, P_k) < nk$, [5], the result is true if $n \le 13k^4$. Thus assume $n > 13k^4$ and the result is true for all appropriate graphs with less than *n* vertices. Let *F* be a graph on $n + 13k^5$ vertices. We will assume that \overline{F} contains no P_k and show that *G* is a subgraph of *F*.

Case I. G has a suspended path with 3k - 3 vertices.

Let H be the graph on n = k + 1 vertices obtained from G by shortening the suspended path by k - 1 vertices. Therefore H has at most $[n(1 + (1/18k))] - k + 1 \le [(n - k + 1)(1 + 1/18k)] + 1$ edges. The induction assumption and Lemma 4 imply that F contains H as a subgraph. Of course G can be obtained from H by lengthening the suspended path in H by k - 1 vertices. Let H' be a subgraph of F in which this suspended path has been lengthened as much as possible (up to k - 1). If H' is isomorphic to G, the proof of this case is complete. If not, then select a set S of k - 1 vertices of F not in H'. Let F' be the subgraph of F spanned by the vertices of the suspended path of H' and the set S. Lemma 1(i) implies that \overline{F}' , and hence \overline{F} , contains a P_k . This contradiction completes the proof of this case.

Case II. G has a vertex adjacent to at least 2k - 2 vertices of degree 1.

Let S be a set of 2k - 2 vertices of degree 1 in G which are adjacent in G to a vertex v. Let u be the endvertex of a path P of maximal length in \overline{F} . The path P has at most k - 1 vertices. If H = G - S, then H has n - 2k + 2 vertices and at most [(n - 2k + 2)(1 + 1/18k)] + 1 edges. The induction assumption and Lemma 4 imply that $r(H, P_k) \le n - k + 1 + 13k^5$. Therefore F contains H as a subgraph such that H is vertex-disjoint from P. Since u is adjacent in F to each vertex of F not on P, replacing v in H by u gives a copy of G in F.

Case III. G has no suspended path with 3k - 3 vertices and no vertex adjacent to 2k - 2 vertices of degree 1.

Lemma 2 implies that G has at least $n/(2(3k-3)) - 3n/36k \ge n/12k$ vertices of degree 1. Since $n > 13k^4$, G has at least $\{\frac{13}{12}k^3\}$ vertices of degree 1. The fact that no vertex of G is adjacent to more than 2k - 3 vertices of degree 1 insures that there are distinct vertices v_1, v_2, \ldots, v_l , $(l \ge 13k^3/(12(2k-3)))$ in G with each v_i adjacent to n_i vertices of degree 1. In fact we can assume that $2k - 2 > n_1$ $\ge n_2 \ge \cdots \ge n_l \ge 1$ and $\{\frac{13}{12}k^3\} = \sum_{i=1}^l n_i$. Let H be the graph obtained from G by deleting those $t = \sum_{i=1}^l n_i$ vertices of degree 1. The graph H has n - t vertices and at most [(n - t)(1 + 1/18k + t/18k)] edges. The induction assumption and Lemma 4 imply that

 $r(H, P_k) \le n - t + 13k^5 + \{t/18k\}(k - 1).$

A straightforward calculation implies that

$$r(H, P_k) \le (n + 13k^5 - ([k/2] - 1)) - (\sum_{i=1}^{\lfloor k/2 \rfloor - 1} n_i)(k - 1).$$

Therefore by Lemma 3,

$$r(G, P_k) \leq n + 13k^5.$$

Proposition 5 will be used heavily in the proof of the following theorem. Also the outlines of the two proofs are similar in that they both use the structure of a sparse connected graph. Some additional notation is needed before the theorem can be stated.

If v is a vertex of a graph G, let G_v be the subgraph of G induced by the vertices distinct from v and not adjacent to v. Let $\alpha'(G) = \min\{\alpha(G_v): v \in V(G)\}$.

THEOREM 6. Let G be a connected graph with n vertices and no more than $[n(1 + 1/81k^5)]$ edges. Then for $k \ge 2$ and $n \ge 352k^{12}$,

$$r(G, P_k) = \max\{n + \lceil k/2 \rceil - 1, n + k - 2 - \alpha'(G) - \delta\},\$$

where $\delta = 0$ if k - 1 divides $n + k - 3 - \alpha'(G)$ and $\delta = 1$ otherwise.

Note that if $\alpha'(G)$ is as large as $\{k/2\}$, then Theorem 6 implies $r(G, P_k) = n + \lfloor k/2 \rfloor - 1$. Also, unless G has a vertex of very large degree, $\alpha'(G) \ge k$. For example, if v is a vertex of degree $< n(1 - 1/81k^5)$, then G_v has $t \ge n/81k^5$ vertices and at most $t + n/81k^5$ edges. Therefore G_v has average degree no more than 4. Thus Turán's Theorem implies that G_v has independence number at least $t/5 \ge n/405k^5 > k$. Thus the only time it will be necessary to calculate $\alpha'(G)$ will be when G has a unique vertex v of very large degree, and in this case $\alpha'(G)$ is just the independence number of G_v .

PROOF (THEOREM 6). The graph $K_{n-1} \cup K_{\lfloor k/2 \rfloor - 1}$ has no connected subgraph with *n* vertices and its complement has no P_k . Let $n + k - 3 - \alpha'(G) - \delta = t(k-1) + s$, where $0 < s \leq k - 1$. Note that

$$n + k - 3 - \alpha'(G) - \delta = (t - k + 2 + s)(k - 1) + (k - 1 - s)(k - 2).$$

Let L be the graph whose complement is the graph

 $(t-k+2+s)K_{k-1} \cup (k-1-s)K_{k-2}$

Clearly \overline{L} contains no P_k . We will show that L does not contain G as a subgraph if $\alpha'(G) \leq k$. Assume G is a subgraph of L and let v be the unique vertex of maximal degree in G. The vertex v is in one of the disjoint independent sets of L, say S. The set S has at least $k - 1 - \delta$ vertices. Any other vertices of G in S must be independent and not adjacent to v. Therefore S contains at most $\alpha'(G) + 1$ vertices of G. Hence L must have at least

$$n + (k - 1 - \delta) - (1 + \alpha'(G)) = n + k - 2 - \alpha'(G) - \delta$$

vertices, a contradiction. The above two examples give the lower bound for $r(G, P_k)$.

The proof of the upper bound for $r(G, P_k)$ will be broken into three cases, just as the proof of Proposition 5. Let F be a graph on

$$\max\{n + \lfloor k/2 \rfloor - 1, n + k - 2 - \alpha'(G) - \delta\}$$

vertices whose complement \overline{F} has no P_k . We will show that G is a subgraph of F. The result is trivial for k = 2, so assume $k \ge 3$.

Case I. G has a suspended path with at least $13k^5 + 3k$ vertices.

Let H be the graph on $n - 13k^5$ vertices obtained from G by shortening the suspended path by $13k^5$ vertices. By Proposition 5, F has H as a subgraph. Let H'

SPARSE GRAPH-PATH OR CYCLE

be a subgraph of F with a maximal number of vertices which can be obtained from H by lengthening the suspended path as much as possible (up to $13k^5$). If H' is G, the proof is complete. If not, Lemma 1(ii), applied to the subgraph of F induced by the vertices of the suspended path of H' along with $\lfloor k/2 \rfloor$ vertices not in H', implies that \overline{F} contains a P_k . This contradiction completes the proof of this case.

Case II. G has a vertex adjacent to at least $13k^5 + k$ vertices of degree 1.

Let S be a set of $13k^5 + k$ vertices of degree 1 adjacent in G to a vertex v and let H = G - S. Thus H has $n - 13k^5 - k$ vertices and $r(H, P_k) \le n - k$ by Proposition 5.

Consider a maximal length path $P = (x_1, x_2, \ldots, x_t)$ in \overline{F} . By assumption t < k. If there exists an i $(1 < i \le t)$ such that x_1x_i and $x_{i-1}x_i$ are edges in \overline{F} , then the vertices of P form a cycle (i.e. $(x_1, x_i, x_{i+1}, \ldots, x_i, x_{i-1}, x_{i-2}, \ldots, x_2, x_1)$ is a C_i). The maximality of the length of P implies that the vertices of P are a component of \overline{F} . If no such i exists, then the sum of the degrees in \overline{F} of x_1 and x_i is $\le t - 1$. This would insure that \overline{F} has a vertex of degree $\le \lfloor k/2 \rfloor - 1$. Therefore, without loss of generality we can assume that \overline{F} is the disjoint union of complete graphs each with less than k vertices or that \overline{F} has a vertex of degree $\le \lfloor k/2 \rfloor - 1$.

Let u be a vertex of \overline{F} of minimum degree, and let N be the vertices of \overline{F} adjacent to u. Then $N \cup \{u\}$ has at most k - 1 vertices. Since $r(H, P_k) \le n - k$, F has a subgraph H which is vertex-disjoint from $N \cup \{u\}$.

If u has degree at most $\lfloor k/2 \rfloor - 1$, then replace the vertex v in H by the vertex u. Since u is adjacent in F to each vertex of H, this gives another copy of H in F. Also, since u is adjacent in F to all but possibly $\lfloor k/2 \rfloor - 1$ vertices, this verifies that G is a subgraph of F.

Assume u has degree at least [k/2]. In this case \overline{F} is the disjoint union of complete graphs each with at most k - 1 vertices. We have that u is in some component C of \overline{F} with $t \le k - 1$ vertices. In fact, if F has $n + k - 2 - \alpha'(G) - \delta$ vertices, then $t \le k - 1 - \delta$. Since $r(H, P_k) \le n - k$, F has a subgraph H which is vertex-disjoint from C. Replace the vertex v of H by u. This gives a copy of H in F, which we will denote by H'. The graph H' has only the vertex u in common with C. Select an independent set B of min $\{\alpha(H'_u), t - 1\}$ vertices in H'_u . Obtain another copy of H, which we will denote by H'', by replacing the vertices B of H' with |B| vertices of C distinct from u. This can be done, since each vertex of C is adjacent in F to each vertex of H' except u. Since $\alpha(H'_u) = \alpha(H_v)$ and $\alpha(G_v) \ge \alpha'(G)$, u is adjacent in F to all except at most $t - 1 - \alpha'(G)$ of the vertices in F - H''. Therefore the vertex u is adjacent in F to at least $13k^5 + k$ vertices not in H''. Hence F contains G as a subgraph.

Case III. G has no suspended path with $13k^5 + 3k$ vertices and no vertex adjacent to $13k^5 + k$ vertices of degree 1.

Lemma 2 implies that G has at least

$$n/(2(13k^5+3k)) - 3n/(2 \cdot 81k^5) \ge (n/k^5)(\frac{1}{27} - \frac{1}{54}) = n/54k^5$$

vertices of degree 1. Since $n \ge 352k^{12}$, G has at least $(\frac{352}{54})k^7$ vertices of degree 1. Thus there exist vertices v_1, v_2, \ldots, v_l in $G(l \ge (\frac{352}{54})k^7/(13k^5 + k))$ such that v_l is

adjacent to n_i vertices of degree 1 with $13k^5 + k > n_1 \ge n_2 \ge \cdots \ge n_l \ge 1$. We can also assume that $\{(\frac{352}{54})k^7\} = \sum_{i=1}^l n_i$.

Let *H* be the graph obtained from *G* by deleting these $t = \sum_{i=1}^{t} n_i$ vertices of degree 1. Therefore *H* has n - t vertices and at most $[(n - t)(1 + 1/81k^5) + (t/81k^5)]$ edges. Thus Proposition 5 and Lemma 4 imply

$$r(H, P_k) \le n - t + 13k^5 + \left[\frac{t}{81k}5\right](k-1).$$

Using very bad approximations, direct calculations verify that

$$r(H, P_k) \leq n - \left(\sum_{i=1}^{\lfloor k/2 \rfloor - 1} n_i\right)(k-1).$$

Therefore by Lemma 3, G is a subgraph of F. This completes the proof of the theorem.

The lower bound on n, the number of vertices in G in Theorem 6, and the upper bound on the number of edges in G in Proposition 5 and Theorem 6 are by no means sharp. Sometimes to simplify the calculations slightly weaker bounds were used, but significant improvement in these bounds will necessitate different methods of proof.

There are several interesting corollaries of Theorem 6. The first one is a consequence of the remarks made prior to the proof of Theorem 6, and the last three are a result of calculating $\alpha'(G)$. The last two are special cases of well-known results.

COROLLARY 7. Let G be a connected graph with n vertices and no more than $[n(1 + 1/81k^5)]$ edges. If $\Delta(G) \leq n(1 - 1/81k^5)$, $n \geq 352k^{12}$ and $k \geq 2$, then

$$r(G, P_k) = n + \lfloor k/2 \rfloor - 1.$$

COROLLARY 8. If T_n is a tree on n vertices with $\Delta(T_n) < n - k$, $n \ge 352k^{12}$ and $k \ge 2$, then

$$r(T_n, P_k) = n + \lceil k/2 \rceil - 1.$$

COROLLARY 9 [6, 7]. If $n \ge 352k^{12}$ and $k \ge 2$, then (i) $r(P_n, P_k) = n + \lfloor k/2 \rfloor - 1$, (ii) $r(C_n, P_k) = n + \lfloor k/2 \rfloor - 1$.

In the results as stated in both [6] and [7] the restrictions on n are only linear in k.

COROLLARY 10 [10]. If $n \ge 352k^{12}$ and $k \ge 2$, then

$$r(S_n, P_k) = n + k - 2 - \delta,$$

where $\delta = 0$ if k - 1 divides n - 2 and $\delta = 1$ otherwise.

In the result stated in [10] the restriction on n is only quadratic in k.

We will use the results on $r(G, P_k)$ to obtain the numbers $r(G, C_k)$ for k odd. Two additional lemmas will be needed. These we will state and prove now.

LEMMA 11. Let G be a graph with n vertices and minimum degree $\delta \ge 1$. If v is a vertex of degree δ , then

$$r(G, C_k) \leq \max(r(G - v, C_k), \delta r(G, P_{k-1}) + n - \delta).$$

PROOF. Let F be a graph with $\max(r(G - v, C_k), \delta r(G, P_{k-1}) + n - \delta)$ vertices whose complement \overline{F} does not contain a C_k . Therefore by assumption, G - v is a subgraph of F. Let N be the vertices of G - v which are adjacent in G to v, and let S be the vertices of F not in G - v. Thus N has δ vertices and S has at least $\delta(r(G, P_{k-1}) - 1) + 1$ vertices. If F does not contain G as a subgraph, then each vertex of S must be adjacent in \overline{F} to at least one vertex of N. Therefore some vertex w of N has degree at least $r(G, P_{k-1})$ in \overline{F} . This implies that either F contains G as a subgraph or there is a C_k in \overline{F} containing w. This completes the proof.

LEMMA 12. If G is a graph with n vertices and l edges, then

$$r(G, C_k) \le n + 2lk - 2l/n.$$

PROOF. The proof will be by induction on *n*. If n = 1, then l = 0 and clearly $r(G, C_k) \leq 1$. Proceed by induction and let *v* be a vertex of minimal degree δ in *G*; Lemma 11 implies

$$r(G, C_k) \leq \max(r(G - v, C_k), \delta \cdot r(G, P_{k-1}) + n - \delta).$$

The induction assumption and the fact that $r(K_n, P_k) \leq nk$ give that

$$r(G, C_k) \leq \max\left(n-1+2k(l-\delta)-\frac{2(l-\delta)}{n-1}, \delta kn+n-\delta\right).$$

Using that $\delta \leq 2l/n$, it is easily verified that $\delta kn + n - \delta \leq n + 2kl - 2l/n$. Therefore direct calculation gives

$$r(G, C_k) \le n + 2kl - \frac{2l}{n},$$

which completes the proof.

Note that Lemma 12 implies for example that $r(K_n, C_k) \le n + 2n(n-1)k/2 \le n^2k$ [1].

PROPOSITION 13. Let G be a connected graph with n vertices and no more than $[n(1 + 1/12k^2)]$ edges. For all $n, k \ge 3$,

$$r(G, C_k) \leq 2n + 13k^5.$$

PROOF. Assume the proposition is not true. Then there exists a graph F on $2n + 13k^5$ vertices which does not have G as a subgraph and whose complement \overline{F} does not contain a C_k . Let L be the graph obtained from G by successively deleting vertices of degree 1 and shortening suspended paths with at least 5[(k + 1)/2] + 1 vertices by a vertex. We will show that L is also not a subgraph of F. This will lead to a contradiction.

Case I. G has a suspended path with at least 5[(k + 1)/2] + 1 vertices.

Let H be the graph obtained from G by shortening the suspended path by a vertex. If F contains H as a subgraph, then Lemma 1 implies that for any pair of

vertices of F not in H there is a path with k (k odd) or k - 1 (k even) vertices between them. This path can be assumed to use any [(k + 1)/2] predetermined vertices of F not in H. Since $r(G, P_k) \le n + 13k^5$ by Proposition 5, either F contains G as a subgraph or \overline{F} contains a C_k using the path described earlier. We can thus conclude that H is not a subgraph of F.

Case II. G has vertex of degree 1.

Let *H* be the graph obtained from *G* by deleting a vertex of degree 1. Let *v* be the vertex of *G* adjacent to this vertex of degree 1. If *F* contains *H* as a subgraph, then the vertex *v* is adjacent in \overline{F} to at least $n + 13k^5 + 1$ vertices. Since $r(G, P_k) \le n + 13k^5$, we can assume that *F* does not contain *H* as a subgraph.

Repeated application of the situations in cases 1 and 2 imply that L is not a subgraph of F. The graph L has l vertices, no more than $l + [n/12k^2]$ edges, no suspended path with more than 5[(k + 1)/2] vertices, and no vertices of degree 1. Lemma 2 implies that l satisfies the inequality

$$l/(2 \cdot 5[(k+1)/2]) - 3n/(2 \cdot 12k^2) \le 0.$$

Therefore $l \le (15n/12k^2)[(k + 1)/2]$. Hence by Lemma 12,

$$r(L, C_k) \leq \frac{15n}{12k^2} \left[\frac{k+1}{2} \right] + 2k \left(\frac{15n}{12k^2} \left[\frac{k+1}{2} \right] + \left[\frac{n}{12k^2} \right] \right) \leq 2n$$

since $k \ge 3$. This gives a contradiction, which completes the proof.

Proposition 13 will be used in the proof of the following theorem. Just as we mentioned earlier, the bounds in Proposition 13 and Theorem 14 are not sharp.

THEOREM 14. Let G be a connected graph with n vertices and no more than $[n(1 + 1/42k^5)]$ edges. Then for odd $k \ge 3$ and $n \ge 756k^{10}$,

$$r(G, C_k) = 2n - 1.$$

PROOF. The graph $K_{n-1} \cup K_{n-1}$ gives that $r(G, C_k) \ge 2n - 1$. Let F be a graph on 2n - 1 vertices whose complement \overline{F} does not contain a C_k . We will show that F contains G as a subgraph.

Case I. G has a suspended path with $[13k^5/2] + 5(k+1)/2 + 1$ vertices.

Let *H* be the graph on $n - [13k^5/2] - 1$ vertices obtained from *G* by shortening the suspended path by $[13k^5/2] + 1$ vertices. Then $r(H, C_k) \le 2n - 1$ by Proposition 13. Therefore *F* has *H* as a subgraph. Let *H'* be the subgraph of *F* which can be obtained from *H* by lengthening the suspended path as much as possible (up to $[13k^5/2] + 1$ vertices). If *H'* is *G*, the proof of this case is complete. If not, then Lemma 1(iii) implies that \overline{F} contains a path with *k* vertices between any pair of vertices not in *H'*. Hence any such pair of vertices must be adjacent in *F*. This gives that *G* is a subgraph of *F* since *F* contains a complete graph on at least *n* vertices.

Case II. G has no suspended path with $[13k^5/2] + 5(k + 1)/2 + 1$ vertices. Lemma 2 implies that G has at least

 $n/(([13k^5/2] + (k + 1)/2 + 1)2) - 3n/((42k^5)2)$

vertices of degree 1. Since $n \ge 756k^{10}$, G has at least $(n/k^5)(\frac{1}{14} - \frac{1}{28}) \ge 27k^5$ vertices of degree 1.

Let *H* be the graph obtained from *G* by deleting $13k^5 + (k-1)/2$ vertices of degree 1. Proposition 13 implies that $r(H, C_k) < 2n - 1$, and thus *H* is a subgraph of *F*. Let *H'* be a maximal connected subgraph of *G* containing *H* which is a subgraph of *F*. If H' = G, the proof is complete, so assume *H'* is a proper subgraph of *G*. Therefore there is a vertex v of *H'* which is adjacent in \overline{F} to each vertex not in *H'*. Let *S* be a subset of V(F) - V(H') with *n* vertices and let R = V(F) - S.

A P_{k-1} in \overline{F} using only the vertices of S would imply that C_k is a subgraph of F. If k = 3, the result is trivial. Assume for the remainder of the proof that $k \ge 5$. Since $r(H, P_{k-1}) \le n - (k-1)/2$ by Proposition 5, there is a copy of H in F using only vertices of S. Denote this copy of H by H_1 . By adding vertices of degree 1 enlarge H_1 as much as possible to obtain a subgraph H'_1 of F with $H_1 \le H'_1 \le G$. If $H'_1 = G$ the proof is complete so assume not. Thus there is a vertex v_1 of H'_1 and of S which is adjacent in \overline{F} to each vertex not in H'_1 . Let N_1 be the vertices adjacent to v_1 in \overline{F} . Hence $|N_1| \ge n$ and $|N_1 \cap R| \ge n - 13k^5 - (k-1)/2$. By induction we will select (k-1)/2 vertices from the set S. Assume v_1, v_2, \ldots, v_i have been selected. Consider $S - \{v_1, v_2, \ldots, v_i\}$ and repeat the procedure for the selection of v_1 using $S - \{v_1, v_2, \ldots, v_i\}$ instead of S. This gives a vertex v_{i+1} . In this way one obtains a set $\{v_1, v_2, \ldots, v_{(k-1)/2}\}$ of vertices of S with corresponding sets N_i $(1 \le i \le (k-1)/2)$, which satisfy $|N_i| \ge n$ and $|N_i \cap R| \ge n - 13k^5 - (k-1)/2$.

Clearly $N = \bigcap_{i=1}^{(k-1)/2} (N_i \cap R)$ has at least

$$n - ((k-1)/2)(13k^{5} + (k-1)/2)$$

vertices. Therefore there is a large complete bipartite graph in \overline{F} since each vertex v_i $(1 \le i \le (k-1)/2)$ is adjacent in \overline{F} to each vertex of N. In fact if $i \ne j$, $u_i \in N_i$, $u_j \in N_j$ and $u_i u_j \in E(\overline{F})$, then \overline{F} contains a C_k using u_i , u_j , and some vertices from N and $\{v_1, v_2, \ldots, v_{(k-1)/2}\}$. Thus we can assume that $u_i u_j \in E(F)$. This implies that F contains a complete graph on the vertices of $N_1 \cap N_2$ and each vertex of N_1 is adjacent in F to each vertex in $N_1 \cap N_2$. Since $|N_1 \cap N_2| > n - 27k^5$ and $|N_1| \ge n$, H is a subgraph of F. This completes the proof.

A tree is a special case of a sparse graph. Therefore we have the following corollary.

COROLLARY 15. If T_n is a tree with n vertices, then for odd $k \ge 3$ and $n \ge 756k^{10}$,

$$r(T_n, C_k) = 2n - 1.$$

It would be nice to know the smallest integer $n_0 = n_0(k)$ such that if $n \ge n_0$ then $r(T_n, C_k) = 2n - 1$.

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