ON THE EXISTENCE OF TWO NON-NEIGHBORING SUBGRAPHS IN A GRAPH

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Does there exist a function f(r, n) such that each graph G with $\chi(G) \ge f(r, n)$ contains either a complete subgraph of order r or else two non-neighboring n-chromatic subgraphs? It is known that f(r, 2) exists and we establish the existence of f(r, 3). We also give some interesting results about graphs which do not contain two independent edges.

1. Introduction

Two subgraphs G_1 , G_2 of a graph G are called non-neighboring if there is no edge $v_1v_2 \in E(G)$ with $v_1 \in G_1$ and $v_2 \in G_2$. In general, an arbitrary graph may not contain two non-neighboring subgraphs at all, for example the complete graph. In this paper we raise the following question: Is there a minimal integer f(r, n) such that each graph G with $\chi(G) \ge f(r, n)$ and which does not contain a complete subgraph of order r must contain two non-neighboring n-chromatic subgraphs? An upper bound for f(r, 2) follows from a result of S. Wagon [2]. Here we show that it is sufficient to prove the existence of f(r, n) for $r \leq n$. More precisely, for a fixed n, an upper bound for f(r, n), r > n, is given in terms of f(r, n), $r \le n$. The proof is based on the same idea of S. Wagon. From $f(3, 3) \leq 8$ we deduce an upper bound for f(r, 3). Next we investigate graphs which do not contain $2K_2$ as an induced subgraph. We say that the two edges v_1v_2 , u_1u_2 of the graph G are independent if the subgraph induced by v_1, v_2, u_1, u_2 is $2K_2$, i.e., the complement of a chordless 4-cycle. We prove that a vertex-critical 4-chromatic graph G which does not contain two independent edges has order $|G| \leq 13$. We also give a lower bound for the maximum degree of a graph without two independent edges.

2. Notation

We consider graphs G = (V(G), E(G)) which are finite, loopless and have no multiple edges. The neighborhood N(v) of a vertex $v \in G$ is the set of vertices adjacent to v. We put $N^*(v) = N(v) \cup \{v\}$. For $W \subseteq V(G)$, we denote $N(W) = \bigcup \{N(v): v \in W\}$ and $N^*(W) = N(W) \cup W$. If G_1 is a subgraph of G then $N(G_1)$, $N^*(G_1)$ respectively denote $N(V(G_1))$, $N^*(V(G_1))$. Two subgraphs G_1 , G_2 are non-neighboring if $V(G_1) \cap$ $\cap N^*(G_2) = \emptyset$. A subset $W \subseteq V(G)$ is called a dominating set if $N^*(W) = V(G)$.

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3. The functions f(r, 2) and f(r, 3)

S. Wagon [2] proved that if G contains neither a complete subgraph of order r nor two independent edges then $\chi(G) \leq {r \choose 2}$. It follows that $f(r, 2) \leq {r \choose 2} + 1$. The slightly stronger result $f(r+1, 2) \leq f(r, 2) + r$ is implicit in [2]. It is trivial that f(2, 2) = 2 and the pentagon C_5 shows that f(3, 2) = 4. The 5-wheel $C_5 + K_1$ shows that $f(4, 2) \geq 5$ and from Wagon's Theorem we have $f(4, 2) \leq 7$. Recently P. Hajnal proved that $f(4, 2) \leq 6$. Finally, Nagy and Szentmiklóssy proved f(4, 2) = 5.

Theorem 1. For r > n,

$$f(r, n) \leq 1 + (n-1)\binom{r-1}{n} + \sum_{j=1}^{n-1} (f(j+1, n)-1)\binom{r-1}{j}.$$

Proof. Let G be a graph which does not contain two non-neighboring *n*-chromatic subgraphs. Let K be a complete subgraph of maximum order in G and assume that

$$|K| = k \ge n. \text{ For each } 1 \le j \le n, \text{ let } S_j^{(i)}, \ 1 \le i \le \binom{k}{j}, \text{ denote the } j\text{-subsets of } V(K). \text{ Put}$$
$$X_n^{(i)} = \{v: N(v) \cap S_n^{(i)} = \emptyset\} \qquad 1 \le i \le \binom{k}{n},$$
$$Y_j^{(i)} = \{v: N(v) \cap V(K) = V(K) - S_j^{(i)}\} \quad 1 \le j < n, \ 1 \le i \le \binom{k}{j}.$$

We have $\chi(X_n^{(i)}) \equiv n-1$ since otherwise, $S_n^{(i)}$ would be non-neighboring to an *n*-chromatic subgraph of $X_n^{(i)}$. Also $\chi(Y_j^{(i)}) \equiv f(j+1,n)-1$ since $Y_j^{(i)}$ does not contain a complete subgraph of order j+1. The union of the $X_n^{(i)}$ and $Y_j^{(i)}$ is V(G). Therefore

$$\chi(G) \leq (n-1)\binom{k}{n} + \sum_{j=1}^{n-1} (f(j+1, n) - 1)\binom{k}{j}$$

which implies the required result.

Theorem 2. $f(3,3) \le 8$.

Proof. Let G be a triangle-free graph which does not contain two non-neighboring odd circuits. Let $C = v_0 v_1 \dots v_{2k}$ be an odd circuit of minimum lenght in G. We describe a proper 7-coloring c of G as follows. Let $c(v_0) = 1$ and $c(v_i) = 2$ (resp. 3) if i is odd (resp. even) and $1 \le i \le 2k$. Further we let c(x) = 2 (resp. 3) if $x \in N(v_i)$ for i even (resp. odd) and $2 \le i \le 2k - 1$. Otherwise, for $x \in N(C)$ we let

$$c(x) = \begin{cases} 1 & x \in N(v_1) \\ 4 & x \in N(v_0) \\ 5 & x \in N(v_{2k}). \end{cases}$$

Since $G - N^*(C)$ does not contain an odd circuit, we need at most two more colors 6 and 7 to extend c to all of V(G). This shows that $\chi(G) \le 7$ and, therefore, $f(3,3) \le 8$.

It is easy to check that the triangle-free 5-chromatic graph described by Mycielski [1] does not contain two non-neighboring odd circuits. This shows that $f(3,3) \ge 6$.

From Theorems 1 and 2 we get the following polynomial upper bound for f(r, 3).

Corollary 3.
$$f(r, 3) \leq 2\binom{r-1}{3} + 7\binom{r-1}{2} + r \quad (r > 3).$$

4. Graphs without two independent edges

In this section we prove some more results about graphs without two independent edges. We start by a result about 4-critical (i.e vertex-critical 4-chromatic) graphs without two such edges. Examples of these graphs are K_4 and the 5-wheel C_5+K_1 . We shall encounter more in what follows. It is somewhat a surprising fact that these graphs cannot have a large order, specially if we know that this is not the case for higher chromatic numbers.

Theorem 4. *If G is a 4-critical graph without two independent edges then* $|G| \leq 13$.

Proof. We may assume that G is not (and therefore does not contain) K_4 . Let $v_1v_2v_3$ be a triangle in G. We have two cases.

Case 1. There is a vertex $v \in G$ adjacent to none of v_1, v_2, v_3 . Since G contains no two independent edges, then each vertex $u \in N(v)$ must be adjacent to exactly two of v_1, v_2, v_3 . Let c be a proper 3-coloring of G-v. There must exist three vertices $u_1, u_2, u_3 \in N(v)$ such that $c(u_i) = i_i (i=1, 2, 3)$. Suppose the vertices v_1, v_2, v_3 were so labelled that $c(v_i) = i$. Thus u is adjacent to v_j iff $i \neq j$. However, G contains no more vertices since, so far, it is 4-critical. G could contain none, one or two more edges connecting some of u_1, u_2, u_3 . These graphs are shown in Figure 1.

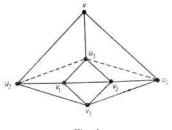


Fig. 1

Case 2. $N^*(v_1, v_2, v_3) = V(G)$. We let

$$A_i = N(v_i) - \bigcup_{j \neq i} N(V_j), \quad B_i = \bigcap_{j \neq i} N(v_j) - \{v_i\}, \quad 1 \le i \le 3,$$
$$A = \bigcup_{i=1}^3 A_i \quad \text{and} \quad B = \bigcup_{i=1}^3 B_i.$$

We note that each B_i is independent since G contains no K_4 . Also, each A_i is independent since if, for example, there were two adjacent vertices $a, a' \in A_1$ then the two edges aa', v_2v_3 would be independent. Therefore both A and B are nonempty. We pick an edge xy with $x \in A_i, y \in B_j$ and $i \neq j$ (if no such edge exists then G is 3-chromatic). For convenience, we assume that $y \in B_3$ and $x \in A_2$. We choose a maximal uniquely 3-colorable subgraph in G - x as follows. Assign color i to v_i (i=1, 2, 3). At each subsequent step a vertex v is assigned color j whenever it is adjacent to two vertices which were previously assigned the two other colors distinct from j. We continue in this way until we cannot proceed further. Denote by W the set of vertices colored in this way and by c(v) the color assigned to $v \in W$. Thus, for example, c(v)=i for each $v \in B_i$. Suppose $c(w) \neq 1$ for each vertex $w \in W$ adjacent to x. Then by putting c(x)=1 and

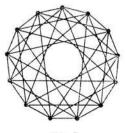
$$c(v) = \begin{cases} 1 & v \in A_2 - W \\ 2 & v \in A_3 - W \\ 3 & v \in A_1 - W \end{cases}$$

we get a proper 3-coloring of G which is a contradiction. Therefore, there is a vertex $w_1 \in W$ adjacent to x with $c(w_1) = 1$. We prove that there is a path $w_1 w_2 \dots w_t$ such that:

- (i) $c(w_i) \equiv i \pmod{3}$
- (ii) for i < t, $w_i \in A_i$ where $i + 2 \equiv j \pmod{3}$ and $w_i \in B_i$ where $t \equiv l \pmod{3}$.

This is true if $w_1 \in B_1$. Suppose not, then necessarily $w_1 \in A_3$ and it was assigned color 1 due to its adjacency to a vertex w_2 (previously) colored with color 2. Either $w_2 \in B_2$ or $w_2 \in A_1$ and we can find w_3 with the required properties. Continuing in this way we, eventually, arrive at $w_t \in B$. Let us assume further that this path is of minimum length. Clearly the vertices $v_1, v_2, v_3, x, y, w_1, \ldots, w_t$ span a 4-chromatic graph so that they must be all of V(G). Thus we have to prove that $t \equiv 8$. Assume, on the contrary, that $t \equiv 9$. Consider the two edges xw_1 and w_7w_8 . There is no edge w_1w_7 since $c(w_1) = c(w_7) = 1$. Also, $xw_7, w_1w_8 \notin E(G)$ since, otherwise, we could have chosen a path of smaller length. Therefore $xw_8 \in E(G)$. Now consider the two edges xw_8 and w_4v_3 . If either $xw_4, w_4w_8 \in E(G)$ then we get a path of smaller length. Also none of x, w_8 is adjacent to v_3 . This is a contradiction since xw_8, w_4v_3 cannot be independent. This completes the proof that $|G| \leq 13$.

To show that 13 in Theorem 4 is best possible, we give a graph G with |G| = 13. This is shown in Figure 2.



In contrast to Theorem 4, we describe a 5-critical graph without two independent edges which has 4n+5 vertices for arbitrary *n*. The vertices of this graph are $x_1, x_2, x_3, x_4, y_0, y_1, \dots, y_{4n}$. The edges are

$$\begin{array}{ll} x_i x_j & (i \neq j), \\ y_i y_{i+1} & (0 \leq i \leq 4n-1), \\ y_i x_j & i-j \equiv 2, 3 \pmod{4}, \\ y_i y_j & i > j \quad \text{and} \quad i-j \equiv 2, 3 \pmod{4}, \\ y_0 x_3 & y_{4n} x_4. \end{array}$$

Our next result describes dominating sets of connected graphs without two independent edges. Here, of course, the connectedness is equivalent to having no isolated vertices.

Theorem 5. Let G be a connected graph without two independent edges. Then G has a dominating set whose induced subgraph is either a complete subgraph or a path on 3 vertices.

Proof. Let $v_1v_2 \in E(G)$. Denote $\lambda = V(G) - N^*(v_1, v_2)$ and $Y = N^*(v_1, v_2) - \{v_1, v_2\}$. We may assume $X \neq \emptyset$ since otherwise $\{v_1, v_2\}$ is a dominating set. The set X is independent but each $x \in X$ is adjacent to at least one $y \in Y$. We choose vertices $y_1, \ldots, \ldots, y_r \in Y$ with r minimum and satisfying:

- (i) $X \subseteq \bigcup_{i=1}^r N(y_i)$,
- (ii) for each *i*, $N(y_i) \cap X$ is maximal that is not properly contained in $N(y) \cap X$ for any $y \in Y$.

If r=1 then $\{v_1, v_2, y_1\}$ is a dominating set with the required property. Let us assume that $r \ge 2$. If $y_i \ne y_j$ then we can find $x, x' \in X$ such that $y_i x, y_j x' \in E(G)$ but $y_i x', y_j x \notin E(G)$. Therefore $y_i y_j \in E(G)$ for all $i \ne j$. Obviously $\{v_1, v_2, y_1, \dots, y_r\}$ is a dominating set. So we need only to prove that for i=1, 2, either v_i is adjacent to all of y_1, \dots, y_r or else $N(v_i) \subseteq N(y_1, \dots, y_r)$. Suppose on the contrary that, for example, $v_1 y_1 \notin E(G)$ and there is a vertex $v \in N(v_1)$ adjacent to none of y_1, \dots, y_r . Let $x \in N(y_1) \cap X$. Since the two edges $v_1 v, y_1 x$ are not independent, then we must have $xv \in E(G)$. Therefore $N(y_1) \cap X \subseteq N(v) \cap X$ and by (ii) above equality holds. Choose two vertices $x_1, x_2 \in X$ with $x_1 y_1, x_2 y_2 \in E(G)$ and $x_1 y_2, x_2 y_1 \notin E(G)$. The two edges $x_1 v, x_2 y_2$ are then independent. This is a contradiction and our theorem is proved.

Corollary 1. If G is a connected graph of order n and without two independent edges then

its maximum degree
$$\Delta(G) \cong \min\left\{2\sqrt{n}-2, \frac{1}{3}(n+1)\right\}.$$

Proof. If G has the vertices of a complete subgraph of order r as a dominating set then this complete subgraph contains a vertex x with degree

$$d(x,G) \cong \frac{1}{r}(n-r)+r-1 \cong 2\sqrt{n}-2.$$

If the dominating set is a path $v_1 v_2 v_3$ then for some *i*

$$d(v_i, G) \ge \frac{1}{3}(n+1).$$

For each *n*, we can construct a graph *G* with |G|=n, $\Delta(G)=\lceil 2\sqrt{n}-2 \rceil$ ([*x*] is the smallest integer $\geq x$) and no two independent edges as follows. Choose two positive integers *r*, *s* with $rs \geq n$ and r+s is minimum. Starting from the vertices v_1, \ldots, v_r of the complete graph K_r , we add n-r more vertices u_1, \ldots, u_{n-r} each joined to at least one v_i in such a way that no v_i is joined to more than s-1 vertices u_j . Clearly for large *n* these are the only extremal graphs to Corollary 1. In particular, for sufficiently large *n*, $\lceil 2\sqrt{n}-2 \rceil$ is the smallest possible maximum degree. However, even for small values of *n*, this is not far from being true.

Corollary 2. All connected graphs G on n vertices and without two independent edges satisfy $\Delta(G) \ge 2\sqrt{n} - 2$ except the three graphs shown in Figure 3.

Proof. Assume $\Delta(G) < 2\sqrt{n} - 2$. Then, since $\Delta(G)$ is an integer, we must have $\left[\frac{1}{3}(n+1)\right] < 2\sqrt{n} - 2$. This is true only for n=5, 7, 8, 10, 11, 13, 14, 17. However,

keeping in mind that G must have a path of three vertices as a dominating set, we can check each case to see that no such graph exists except when n=5, 7, 10 where there is a unique graph in each case.



Fig. 3

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